TMA947/MMG621 NONLINEAR OPTIMISATION

Date: 18-08-21 **Time:** $8^{30}-13^{30}$

Aids: Text memory-less calculator, English–Swedish dictionary

Number of questions: 7; passed on one question requires 2 points of 3.

Questions are not numbered by difficulty.

To pass requires 10 points and three passed questions.

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Result announced: 18–09–09

Short answers are also given at the end of the exam on the notice board for optimization

in the MV building.

Exam instructions

When you answer the questions

Use generally valid theory and methods. State your methodology carefully.

Only write on one page of each sheet. Do not use a red pen.

Do not answer more than one question per page.

At the end of the exam

Sort your solutions by the order of the questions.

Mark on the cover the questions you have answered.

Count the number of sheets you hand in and fill in the number on the cover.

Question 1

(the simplex method)

Consider the following linear program:

maximize
$$z = 3x_1 + 2x_2$$
,
subject to $2x_1 + 3x_2 \le 1$,
 $x_1 - x_2 \ge 4$,
 $x_1, x_2 \ge 0$.

(2p) a) Solve the problem using phase I and phase II of the simplex method.

Aid: You may utilize the identity

$$\left(\begin{array}{cc} a & b \\ c & d \end{array}\right)^{-1} = \frac{1}{ad-bc} \left(\begin{array}{cc} d & -b \\ -c & a \end{array}\right).$$

(1p) b) If an optimal solution exists, then use your calculations to decide whether it is unique or not. If the problem is unbounded, then use your calculations to specify a direction of unboundedness of the objective value.

Question 2

Suppose we have an LP problem (the primal one)

minimize
$$c^{T}x$$
,
subject to $Ax \ge b$,
 $x \ge 0^{n}$.

- (1p) State the LP dual problem.
- (2p) Suppose that this dual problem is feasible. Prove, or disprove, whether the primal problem has an optimal solution, or not. Motivate clearly.

Question 3

(feasible direction methods)

Consider the problem to

minimize
$$f(\mathbf{x}) := (x_1 - 1/2)^2 + x_2^2$$
,
subject to $0 \le x_1 \le 1$,
 $0 \le x_2 \le 1$.

- (2p) Draw the first two iterates obtained using the Frank-Wolfe algorithm starting in $x_0 = (0, 1)^{\mathrm{T}}$.
- (1p) Draw the first two iterates obtained using the simplicial decomposition algorithm starting in $\mathbf{x}_0 = (0, 1)^{\mathrm{T}}$.

(on the SQP algorithm and the KKT conditions)

Consider the following nonlinear programming problem: find $x^* \in \mathbb{R}^n$ that solves the problem to

minimize
$$f(\mathbf{x})$$
, (1a)

subject to
$$g_i(\boldsymbol{x}) \le 0, \qquad i = 1, \dots, m,$$
 (1b)

$$h_j(\boldsymbol{x}) = 0, \qquad j = 1, \dots, \ell,$$
 (1c)

where $f: \mathbb{R}^n \to \mathbb{R}$, g_i , and $h_j: \mathbb{R}^n \to \mathbb{R}$ are given functions in C^1 on \mathbb{R}^n .

We are interested in establishing that the classic SQP subproblem tells us whether an iterate $x_k \in \mathbb{R}^n$ satisfies the KKT conditions or not, thereby establishing a natural termination criterion for the SQP algorithm.

Given the iterate x_k , the SQP subproblem is to

$$\underset{p}{\text{minimize}} \ \frac{1}{2} \boldsymbol{p}^{\mathrm{T}} \boldsymbol{B}_{k} \boldsymbol{p} + \nabla f(\boldsymbol{x}_{k})^{\mathrm{T}} \boldsymbol{p}, \tag{2a}$$

subject to
$$g_i(\boldsymbol{x}_k) + \nabla g_i(\boldsymbol{x}_k)^{\mathrm{T}} \boldsymbol{p} \leq 0, \qquad i = 1, \dots, m,$$
 (2b)
 $h_j(\boldsymbol{x}_k) + \nabla h_j(\boldsymbol{x}_k)^{\mathrm{T}} \boldsymbol{p} = 0, \qquad j = 1, \dots, \ell,$ (2c)

$$h_j(\boldsymbol{x}_k) + \nabla h_j(\boldsymbol{x}_k)^{\mathrm{T}} \boldsymbol{p} = 0, \qquad j = 1, \dots, \ell,$$
 (2c)

where the matrix $\boldsymbol{B}_k \in \mathbb{R}^{n \times n}$ is symmetric and positive semi-definite.

Establish the following statement: the vector \boldsymbol{x}_k is a KKT point in the problem (1) if and only if $p = 0^n$ is a globally optimal solution to the SQP subproblem (2). In other words, the SQP algorithm terminates if and only if x_k is a KKT point.

Hint: Compare the KKT conditions of (1) and (2).

(modelling)

A small municipality is forced to close one or several schools. Out of ten existing schools, at most three schools can be closed. The annual cost to keep school i open is c_i kr, where i = 1, 2, ..., 10. School i can educate a maximum of k_i students.

The municipality is divided into J home areas and there is a requirement that all students in an area belong to the same school. There are b_j students in area j and the average distance from area j to school i is d_{ij} km. The estimated annual cost for student travels is set to m kr per km and student.

Formulate a *linear integer program* to decide on which schools to keep and which ones to close, such that we minimize the total cost for schools and travels and fulfill the above listed requirement

Question 6

(true or false)

Indicate for each of the following three statements whether it is true or false. Motivate your answers!

- (1p) a) For the phase I (when a BFS is *not* known a priori) problem of the simplex algorithm, the optimal value is always zero.
- (1p) b) Suppose that the function $f: \mathbb{R}^n \to \mathbb{R}$ is continuously differentiable on \mathbb{R}^n and let G be a symmetric and positive definite matrix of dimension $n \times n$. Then, if $\nabla f(x) \neq \mathbf{0}^n$ and the vector \mathbf{d} fulfils $G\mathbf{d} = -\nabla f(x)$ it holds that $f(x + t\mathbf{d}) < f(x)$ for small enough values of t > 0.
- (1p) c) If the function $g: \mathbb{R}^n \to \mathbb{R}$ is concave on \mathbb{R}^n and $c \in \mathbb{R}$, then the set $\{ \boldsymbol{x} \in \mathbb{R}^n \mid g(\boldsymbol{x}) \leq c \}$ is convex.

(the gradient projection algorithm)

The gradient projection algorithm is a generalization of the steepest descent method to problems over convex sets. Given a feasible point \mathbf{x}_k , the next point is obtained according to $\mathbf{x}_{k+1} = \operatorname{Proj}_X(\mathbf{x}_k - \alpha_k \nabla f(\mathbf{x}_k))$, where X is the convex set over which we minimize, $\alpha_k > 0$ is the step length, and $\operatorname{Proj}_X(\mathbf{y}) = \arg \min_{\mathbf{x} \in X} ||\mathbf{x} - \mathbf{y}||$ denotes the closest point to \mathbf{x} in X.

Consider the problem to

minimize
$$f(\boldsymbol{x}) := x_1^2 + 2x_2^2 - 2x_1x_2 - 2x_1 - 3x_2$$
, subject to $0 \le x_1 \le 3$, $0 < x_2 < 2$.

Start at the point $\mathbf{x}_0 = (0,0)^{\mathrm{T}}$ and perform two iterations of the gradient projection algorithm using step length $\alpha_k = 1$ for all k. You may solve the projection problem in the algorithm graphically. Is the point obtained a global/local minimum? Motivate why/why not.

Chalmers/GU Mathematics

EXAM SOLUTION

TMA947/MMG621 NONLINEAR OPTIMISATION

Date: 18–08–21

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Note: the solutions presented here are brief in relation to the requirements on your answers, in particular regarding your motivations.

(the simplex method)

Rewrite the problem into standard form by and adding/subtracting slack variables s_1 and s_2 to the left-hand side in the first and second constraint, respectively. Moreover, let z := -z to get the problem on minimization form. Thus, we get the following linear program:

minimize
$$z = -3x_1 - 2x_2$$
,
subject to $2x_1 + 3x_2 + s_1 = 1$,
 $x_1 - x_2 - s_2 = 4$,
 $x_1, x_2, s_1, s_2 \ge 0$.

Introducing the artificial variable a, phase I gives the problem

minimize
$$w = a$$
,
subject to $2x_1 + 3x_2 + s_1 = 1$,
 $x_1 - x_2 - s_2 + a = 4$,
 $x_1, x_2, s_1, s_2, a \ge 0$.

Using the starting basis $(s_1, a)^T$ gives

$$m{B} = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}, m{N} = \begin{pmatrix} 2 & 3 & 0 \\ 1 & -1 & -1 \end{pmatrix}, m{x}_B = \begin{pmatrix} 1 \\ 4 \end{pmatrix}, m{c}_B = \begin{pmatrix} 0 \\ 1 \end{pmatrix}, m{c}_N = \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix}.$$

The reduced costs, $\bar{\boldsymbol{c}}_N^T = \boldsymbol{c}_N^T - \boldsymbol{c}_B^T \boldsymbol{B}^{-1} \boldsymbol{N}$, for this basis is $\bar{\boldsymbol{c}}_N^T = \begin{pmatrix} -1 & 1 & 1 \end{pmatrix}$, which means that x_1 enters the basis. The minimum ratio test implies that s_1 leaves.

Updating the basis we now have $(x_1, a)^T$ in the basis. Calculating the reduced costs, we obtain $\bar{\mathbf{c}}_N^T = \begin{pmatrix} 5/2 & 1/2 & 1 \end{pmatrix}$, meaning that the current basis is optimal. The optimal solution is thus $a^* = 7/2$, since $a^* > 0$, so the original problem is infeasble.

Since the original problem is infeasible, so it is neither existing an optimal solution nor unbounded.

Question 2

(1p) a) The LP dual problem is to:

miximize
$$\boldsymbol{b}^{\mathrm{T}}\boldsymbol{y}$$
,
subject to $A^{\mathrm{T}}\boldsymbol{y} \leq c$,
 $\boldsymbol{y} > 0^{m}$.

(2p) b) If the dual problem has a finite optimal solution, then so does the primal problem. If the dual problem is unbounded, then the primal problem is infeasible. See Theorem 10.6 (Strong Duality Theorem).

Question 3

(feasible direction methods)

- (2p) a) For the Frank-Wolfe algorithm, $y_1 = (1,0)^T$, $x_1 = (0,1)^T$, $y_2 = (0,0)^T$, $x_2 = (9/20, 3/20)^T$.
- (1p) b) For the simplicial decomposition algorithm, $P_0 = \emptyset$, $y_1 = (1,0)^T$, $P_1 = (1,0)^T$, $x_1 = (3/4,1/4)^T$, $y_2 = (0,0)^T$, $P_2 = (1,0)^T \bigcup (0,0)^T$, $x_2 = (1/2,0)^T$,

(on the SQP algorithm and the KKT conditions)

The result is based on a comparison between the KKT conditions of the original problem,

minimize
$$f(\mathbf{x})$$
, (1a)

subject to
$$g_i(\boldsymbol{x}) \le 0, \qquad i = 1, \dots, m,$$
 (1b)

$$h_j(\boldsymbol{x}) = 0, \qquad j = 1, \dots, \ell,$$
 (1c)

and those of the SQP subproblem,

minimize
$$\frac{1}{2} \boldsymbol{p}^{\mathrm{T}} \boldsymbol{B}_k \boldsymbol{p} + \nabla f(\boldsymbol{x}_k)^{\mathrm{T}} \boldsymbol{p},$$
 (2a)

subject to
$$g_i(\boldsymbol{x}_k) + \nabla g_i(\boldsymbol{x}_k)^{\mathrm{T}} \boldsymbol{p} \leq 0, \qquad i = 1, \dots, m,$$
 (2b)

$$h_j(\boldsymbol{x}_k) + \nabla h_j(\boldsymbol{x}_k)^{\mathrm{T}} \boldsymbol{p} = 0, \qquad j = 1, \dots, \ell.$$
 (2c)

We first note that the latter problem is a convex one (the matrix \boldsymbol{B}_k was assumed to be positive semidefinite), and that the solution \boldsymbol{p}_k is characterized by its KKT conditions, since the constraints are linear (so that Abadie's CQ is fulfilled). It remains to compare the two problems' KKT conditions. With $\boldsymbol{p}_k = \boldsymbol{0}^n$ they are in fact identical!

(modelling)

Sets:

 $I := \{1, ..., 10\}$, the set of schools.

The decision variables are:

$$x_i = \begin{cases} 1 & \text{keep school } i, \\ 0 & \text{otherwise,} \end{cases}$$

where $i \in I$.

$$y_{ij} = \begin{cases} 1 & \text{home area } j \text{ go to school } i, \\ 0 & \text{otherwise,} \end{cases}$$

where $i \in I$, $j \in J$.

Model:

$$\begin{aligned} & \text{minimize} & & x_i c_i + m b_j d_{ij}, \\ & \text{subject to} & & \sum_{i \in I} x_i \leq 9, \\ & & \sum_{i \in I} x_i \geq 7, \\ & & \sum_{j \in J} b_j y_{i,j} \leq k_i, & i \in I, \\ & & \sum_{j \in J} b_j y_{i,j} \leq k_i, & i \in I, \\ & & \sum_{j \in I} y_{i,j} \leq x_i, & i \in I, j \in J, \\ & & \sum_{i \in I} y_{i,j} = 1, & j \in J, \\ & & x_i \in \{0,1\}, & i \in I, \\ & & d_{ij} \in \{0,1\}, & i \in I, j \in J. \end{aligned}$$

Question 6

(true or false)

- (1p) a) False. The original problem can be infeasible, which means the optimal value for phase I is higher than 0, like question 1 in this exam.
- (1p) b) True. Since $\nabla f(\boldsymbol{x}) \neq \mathbf{0}^n$, and \boldsymbol{G} is a symmetric and positive definite matrix of dimension $n \times n$, we have that $\nabla f(\boldsymbol{x})^{\mathrm{T}} \boldsymbol{d} = -\nabla f(\boldsymbol{x})^{\mathrm{T}} \boldsymbol{G}^{-1} \nabla f(\boldsymbol{x}) < 0$, so \boldsymbol{d} is a decent direction. By defination of decent direction, the clam is true.
- (1p) c) False. For example, $g(\mathbf{x}) = -x^2$ is concave, but $\{\mathbf{x} \in \mathbb{R}^n \mid g(\mathbf{x}) \le -1\}$ is not convex.

(the gradient projection algorithm)

$$\nabla f(\boldsymbol{x}) = (2x_1 - 2x_2 - 2, 4x_2 - 2x_1 - 3)^{\mathrm{T}}, x_0 - \alpha_k \nabla f(\boldsymbol{x}_0) = (2, 3)^{\mathrm{T}}, x_1 = (2, 2)^{\mathrm{T}}, x_1 - \alpha_k \nabla f(\boldsymbol{x}_1) = (4, 1)^{\mathrm{T}}, x_2 = (3, 1)^{\mathrm{T}}.$$

Since the feasible set is convex, there exists an interior point, so the Slater CQ holds. Since it is a convex problem, so the KKT conditions are both necessary and sufficient. $(3,1)^{T}$ is not a KKT point, so it is neither a global nor a local minimum.